

Achieving reliable high performance in LFNs (long-fat networks)

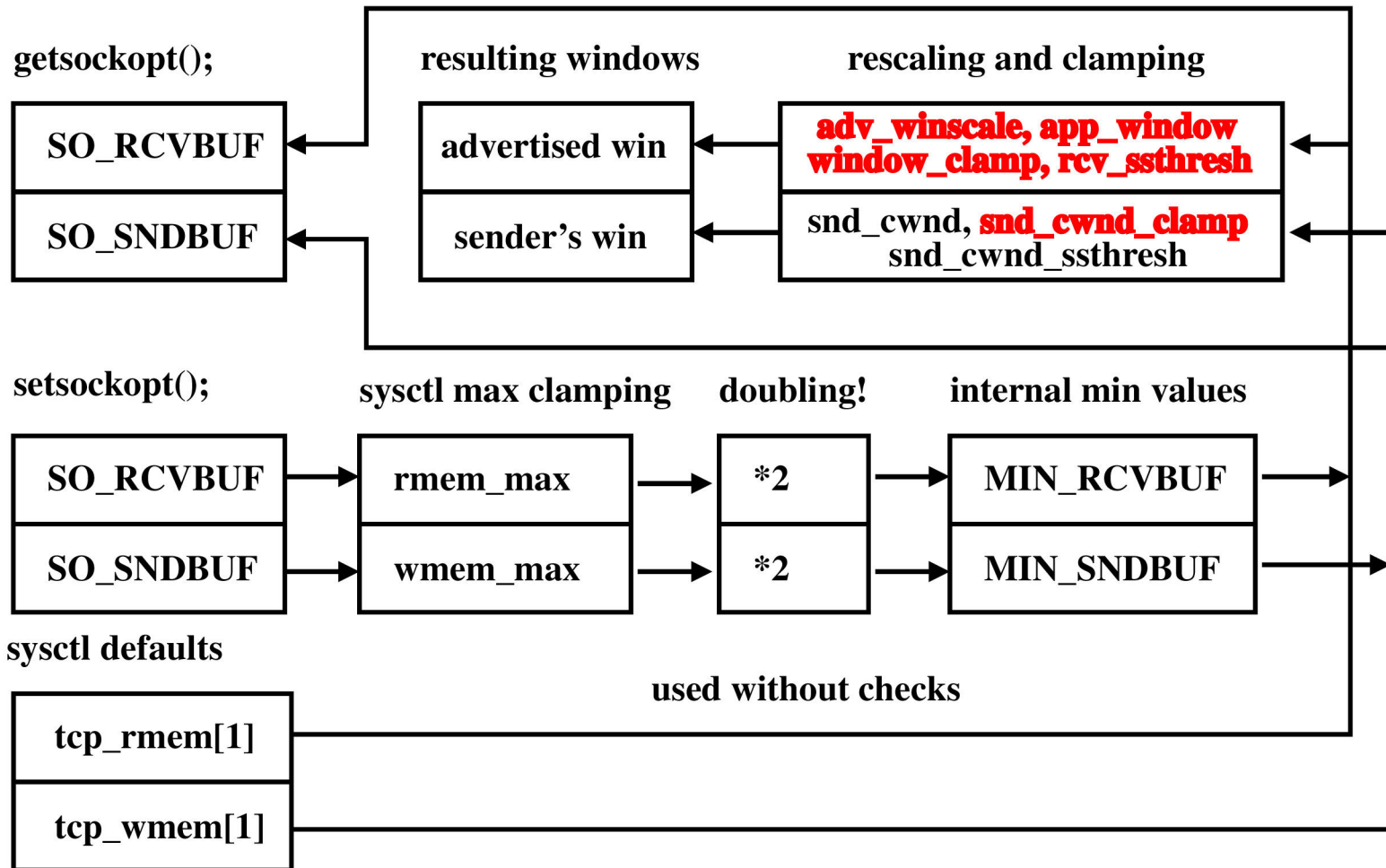
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CESNET

End-to-end performance

- E2E performance is a result of interaction of all computer system components:
 - network, network adapter, communication protocols, operating system, applications
- Decided to concentrate on E2E performance on Linux PCs
- Primary problem areas:
 - TCP window configuration
 - OS / network adapter interaction
 - ssh / TCP interaction

TCP windows configuration

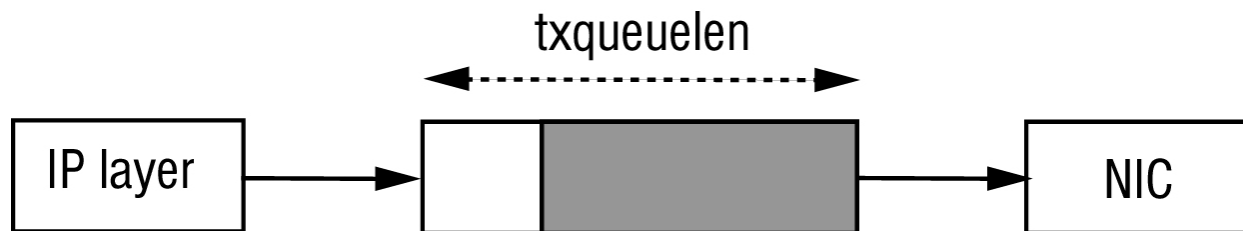
Linux 2.4: how big are my TCP windows?



Throughput with large TCP windows

Interaction with network adapter must be considered:

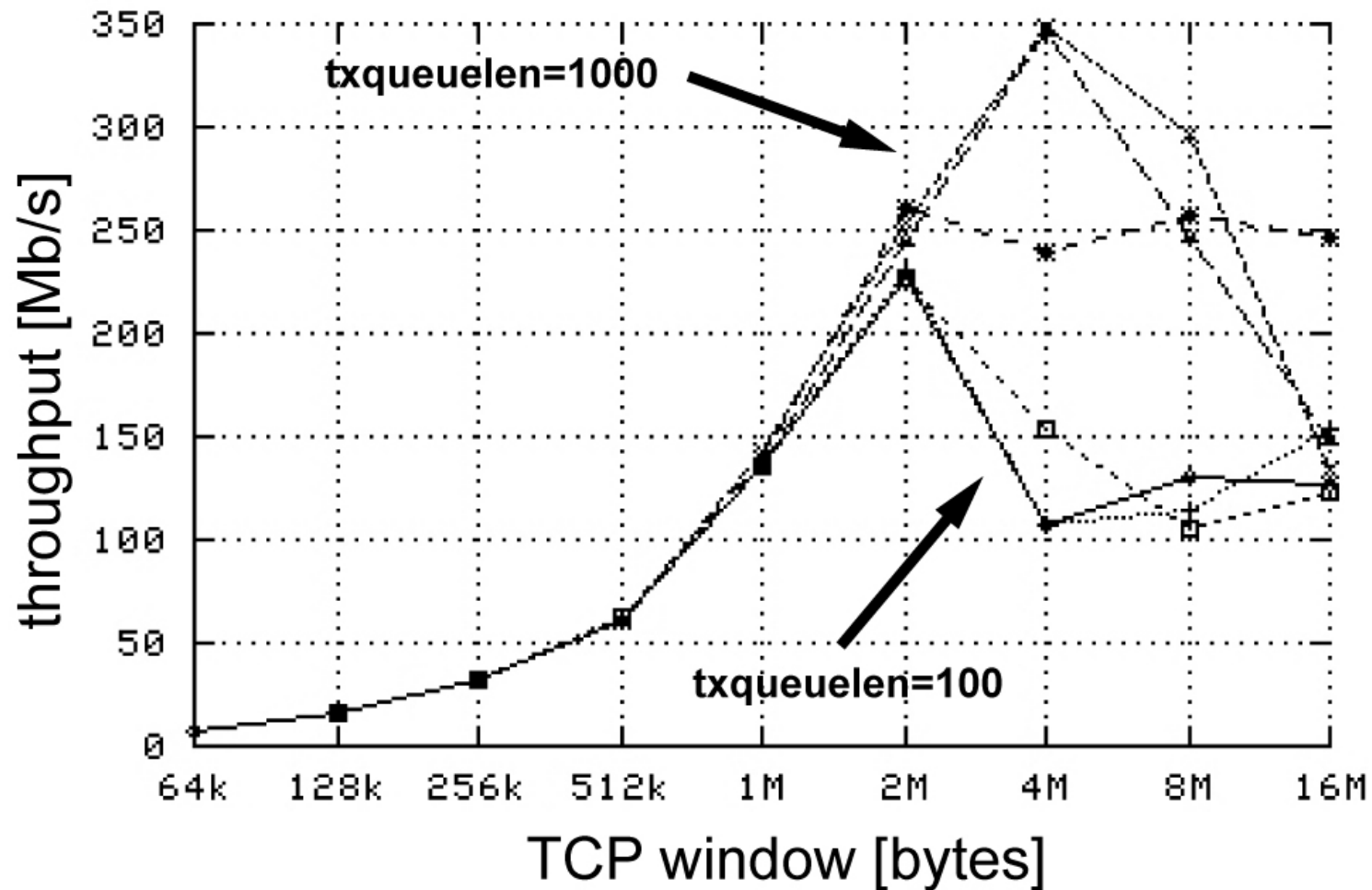
- large TCP sender window allows large chunks of data submitted from IP through txqueue to adapter
- full txqueue -> send_stall() to application and context switch
- no problem as long as txqueue is large enough for a timeslice



for Gigabit Ethernet adapter and standard Linux system timer:
 $\text{txqueuelen} > 1 \text{ Gb/s} * 10 \text{ ms} / 8 \text{ bits} / 1500 \text{ bytes} = 833 \text{ packets}$

```
ifconfig eth0 txqueuelen 1000
```

Throughput with large TCP windows, cont.



Using “buffered pipe” is not good

Router queues must be considered:

- No increase in throughput over using „wire pipe“
- Self-clocking adjusts sender to bottleneck speed, but does not stop sender from accumulating data in queues
- Filled-up queues are sensitive to losses caused by cross-traffic
- Check throughput (TCP Vegas) or RTT increase ?

$rwnd \leq \text{pipe capacity}$

$bw = rwnd / rtt$

$rwnd > \text{pipe capacity}$

$bw \sim (mss / rtt) * 1 / \sqrt{p}$

Flat lower bound

RTT=45ms

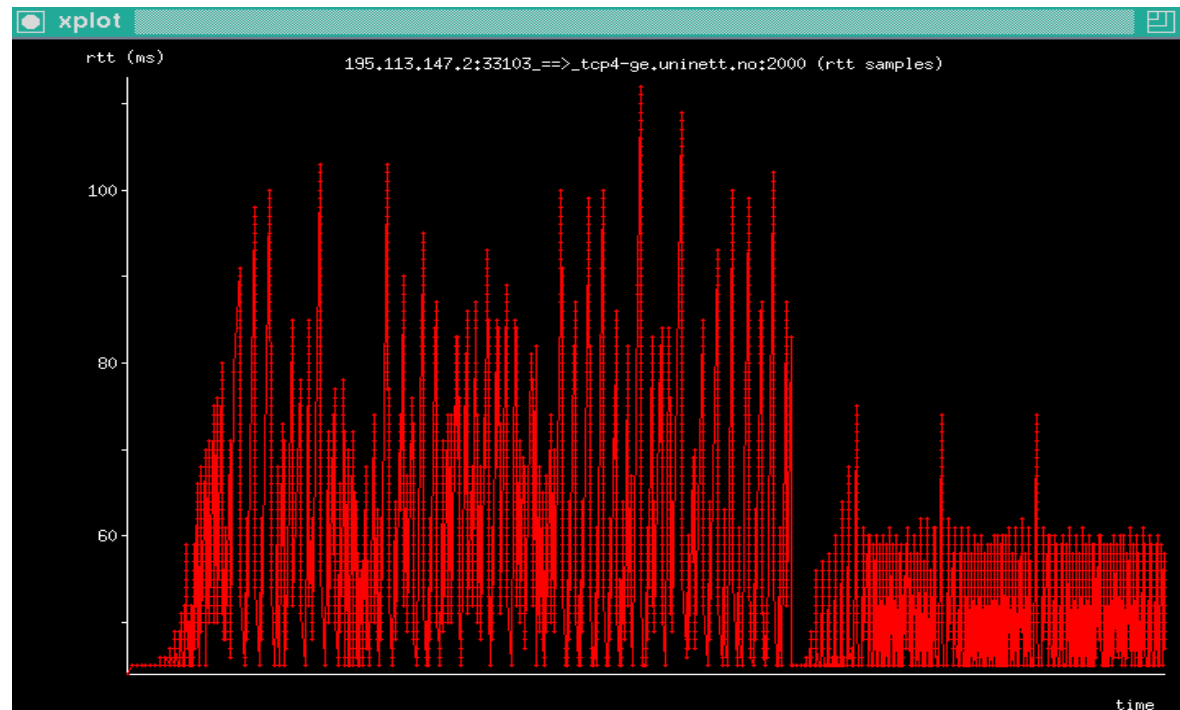
Fluctuations up to

RTT=110ms

Bottleneck

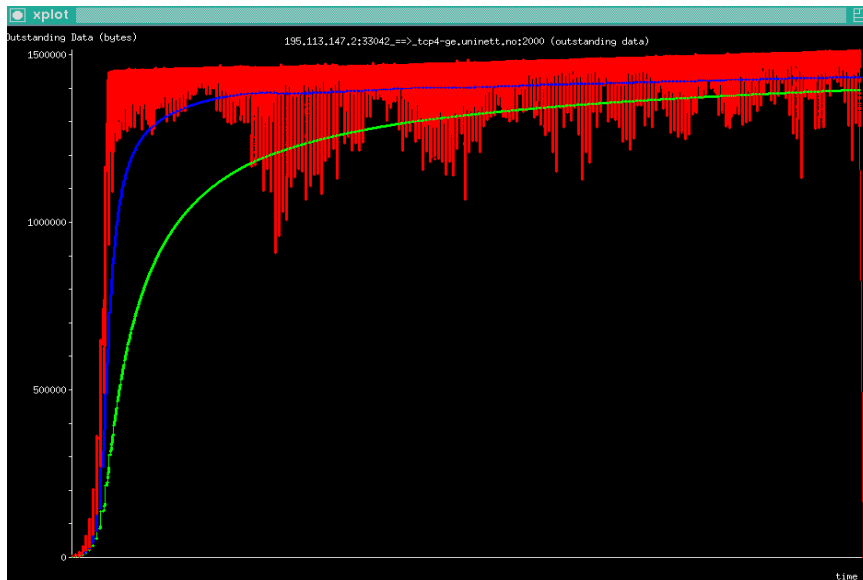
installed BW=1 Gb/s

Buffer content ~8 MB

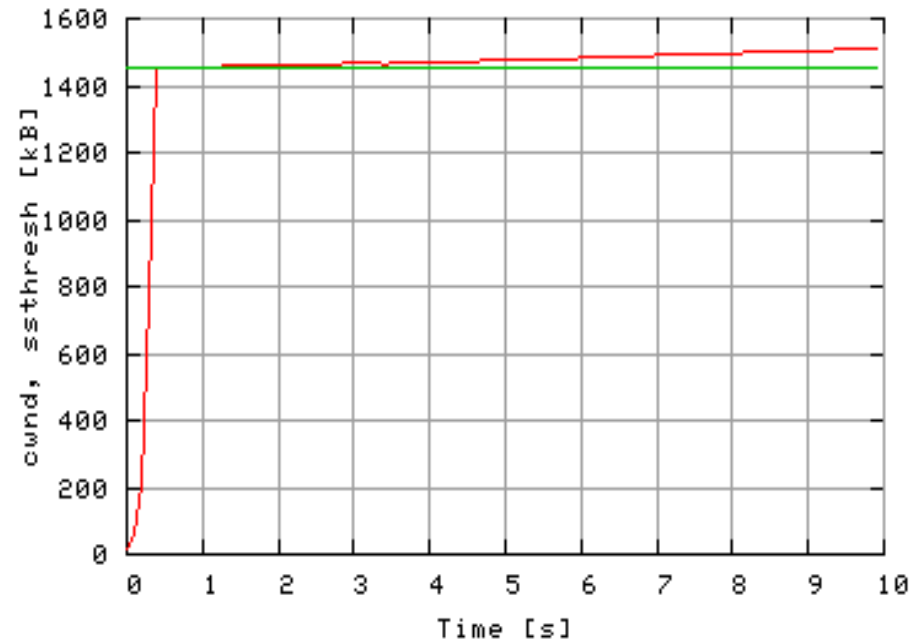


Other configuration problems

TCP cache must be considered



owin development



rwin development

initial ssthresh locked at 1.45 MB

```
echo 1 > /proc/sys/net/ipv4/route/flush
```

Bandwidth measurement and estimation

Test paths: cesnet.cz <--> uninett.no, switch.ch

pathload over one week:

- 27% measurements too low (50-70 Mb/s)
- 7% measurements too high (1000 Mb/s)
- 66% measurements realistic (750-850 Mb/s),
but range sometimes too wide (150 Mb/s)

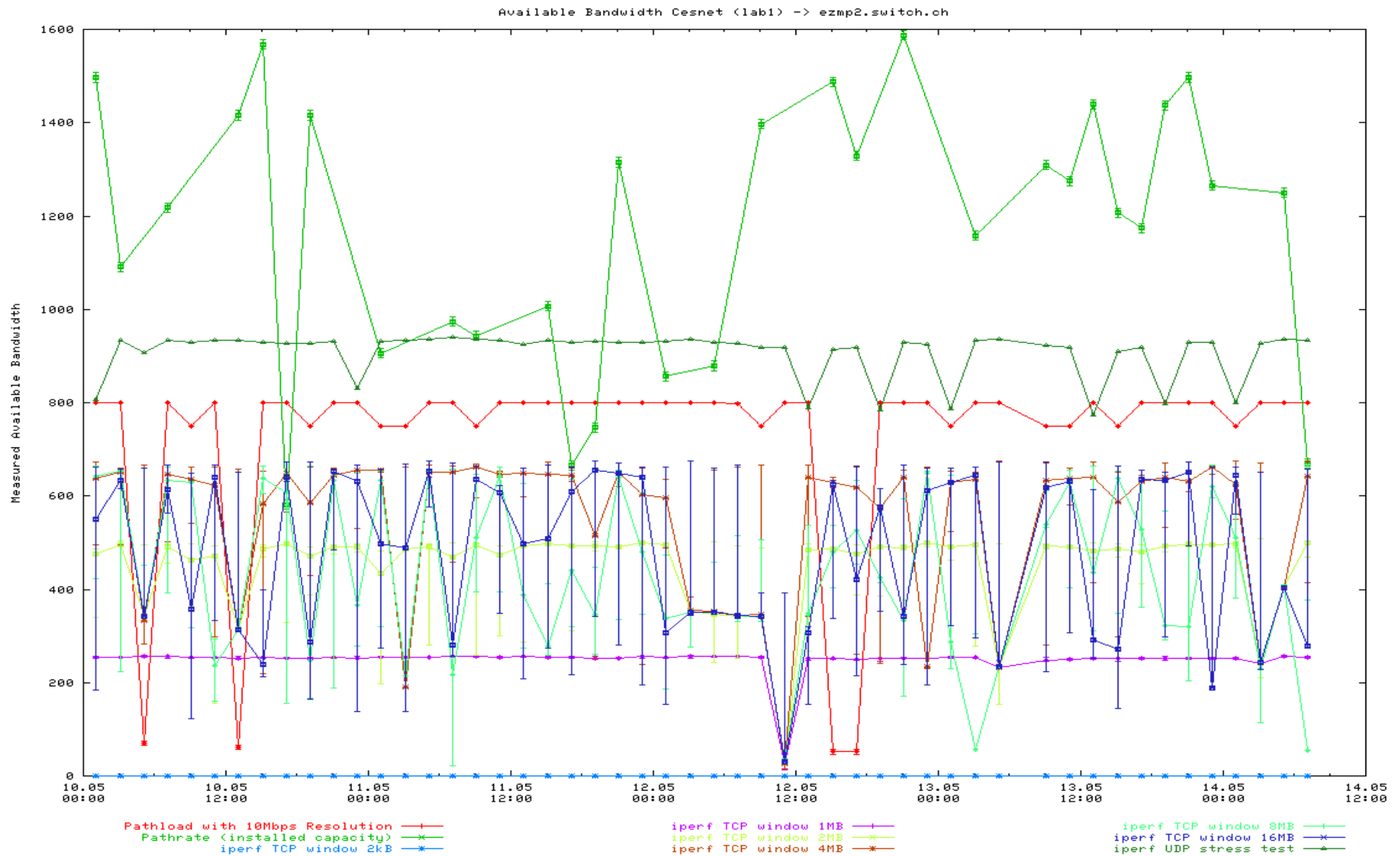
pathrate: lots of fluctuations

UDP iperf: can stress existing traffic

TCP iperf: [more fluctuations for larger TCP windows](#)

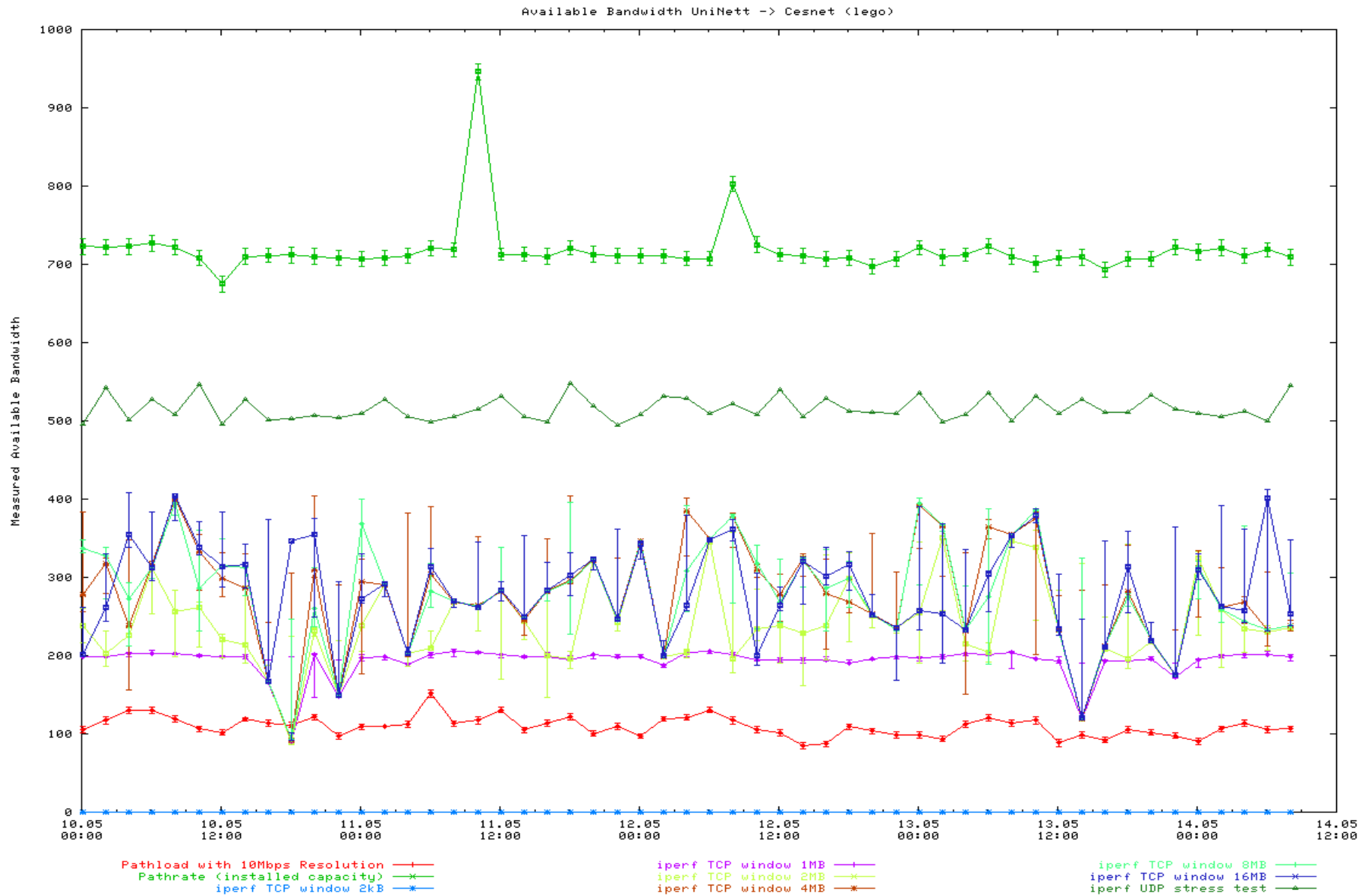
Bandwidth measurement and estimation, cont.

cesnet.cz -> switch.ch



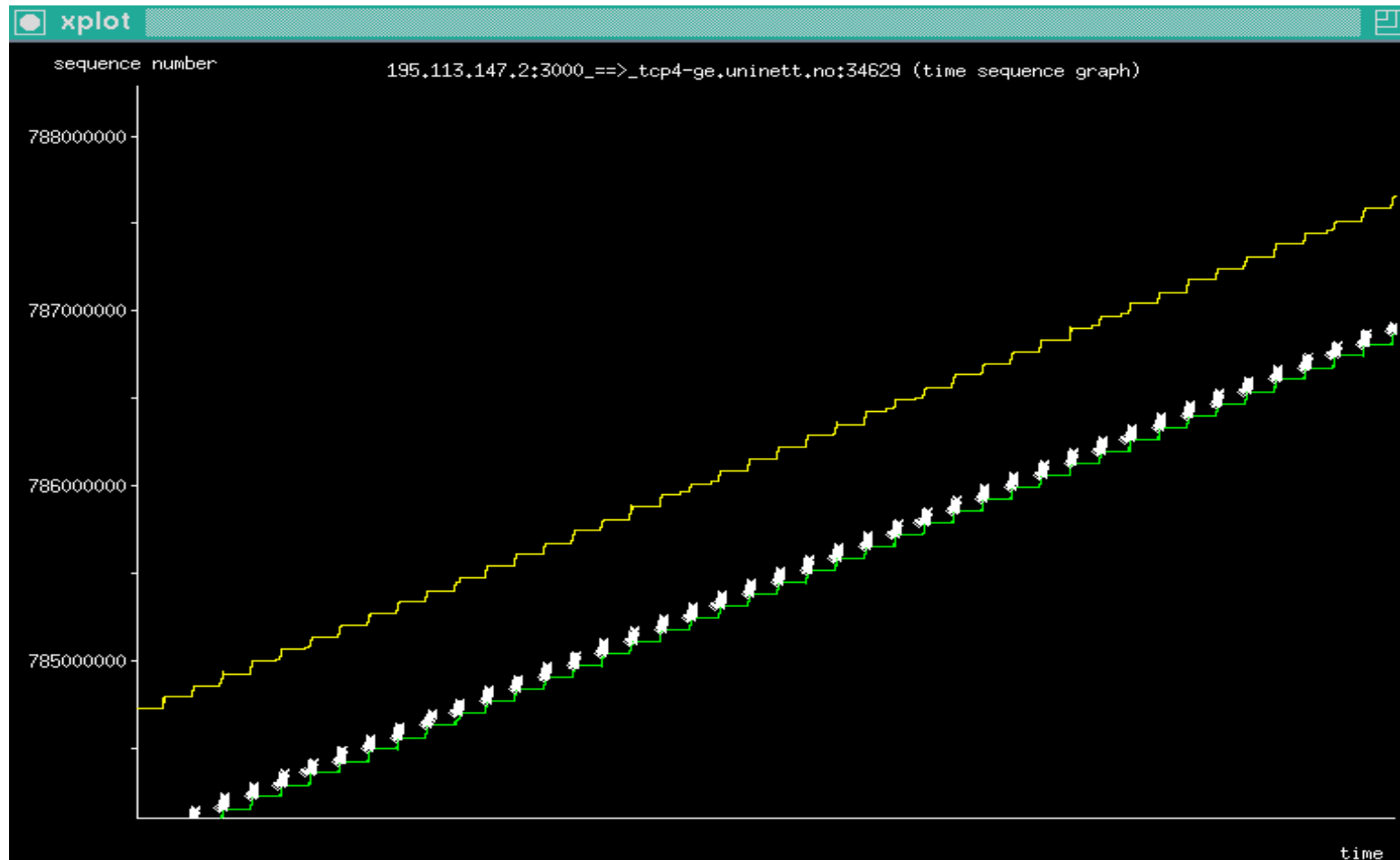
Bandwidth measurement and estimation, cont.

uninett.no -> cesnet.cz



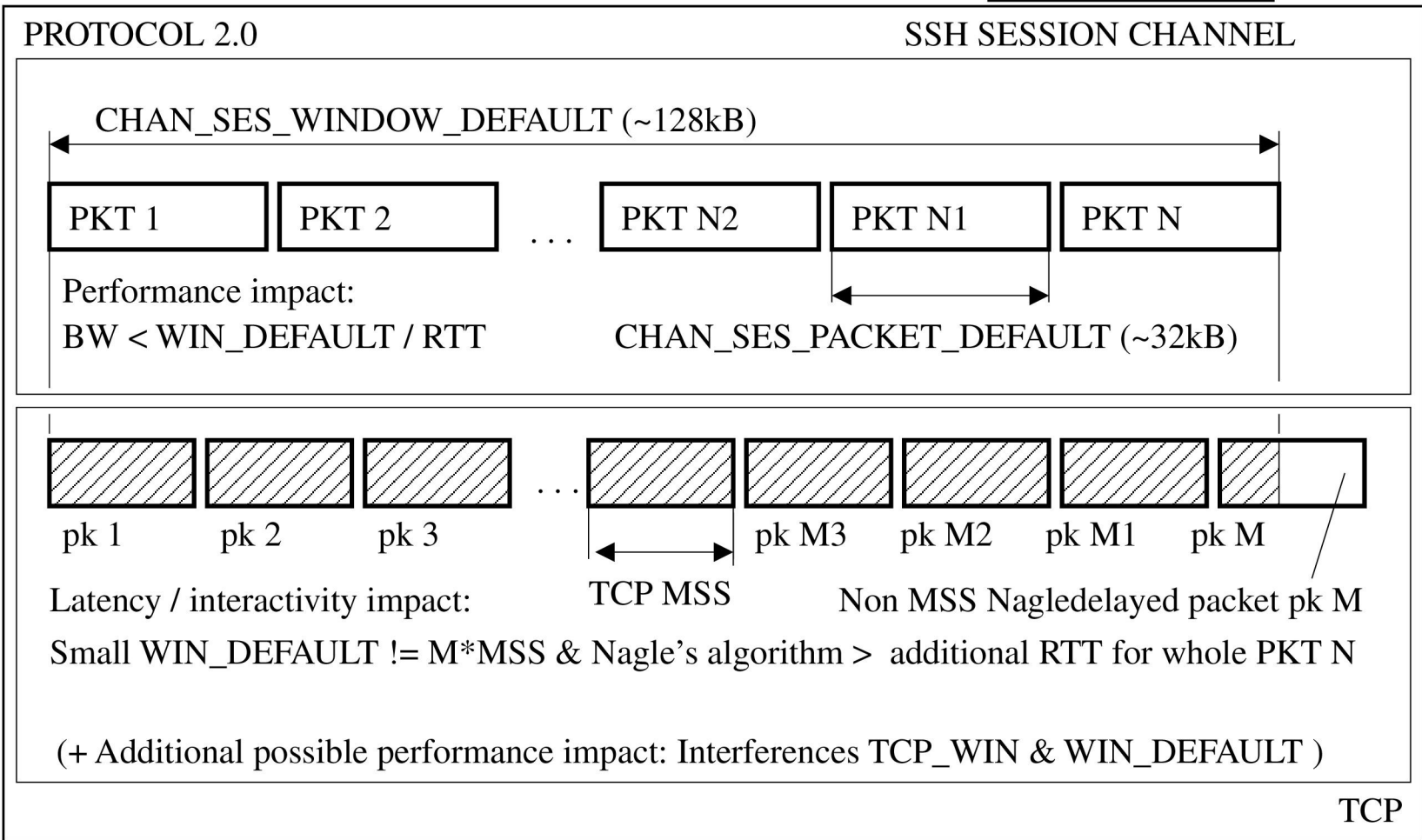
ssh performance

Cesnet -> Uninett, 1.5 MB window, 10.4 Mb/s, 9% load CPU



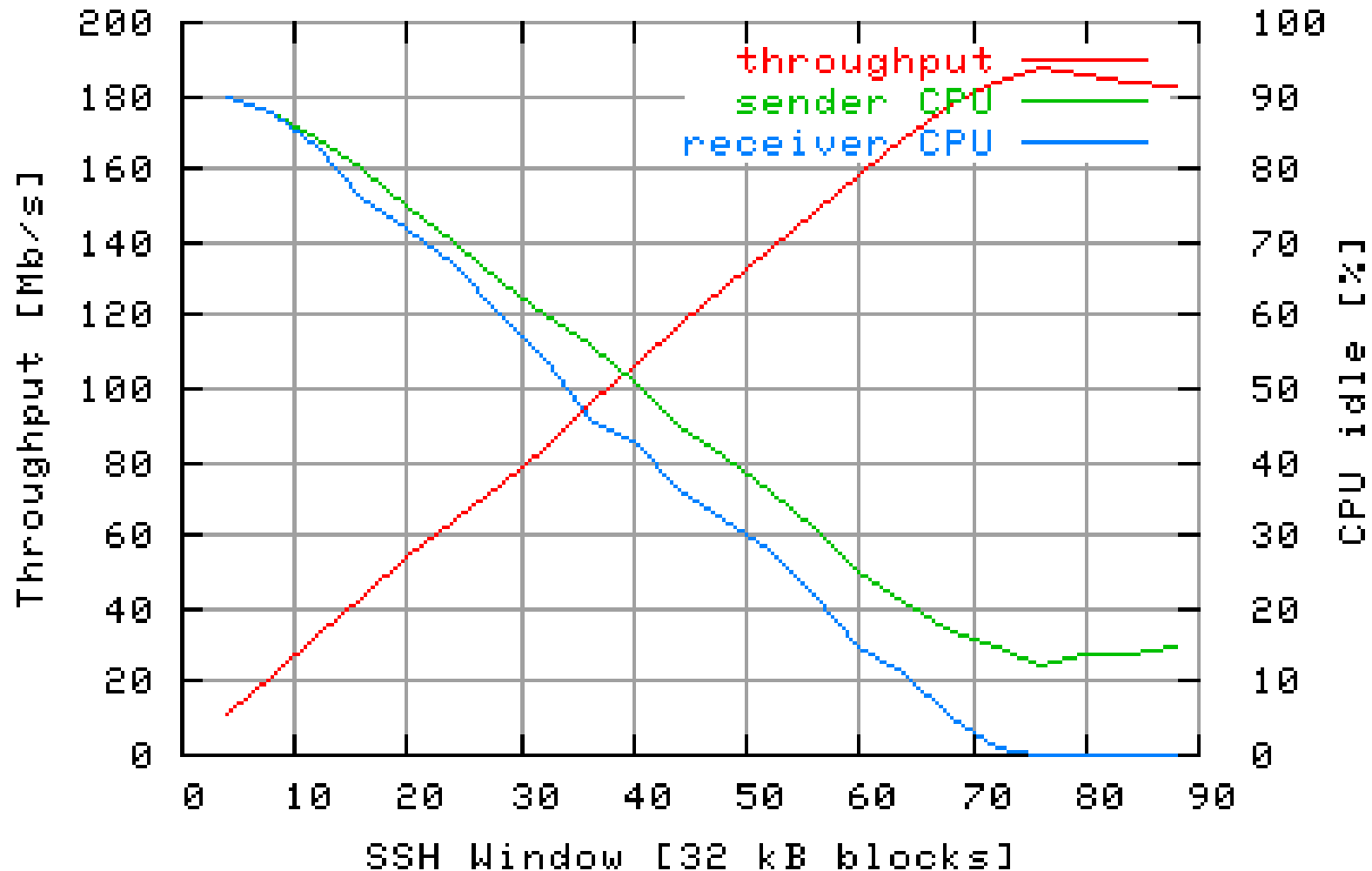
sequence number development - rwin not utilised

ssh performance, cont.



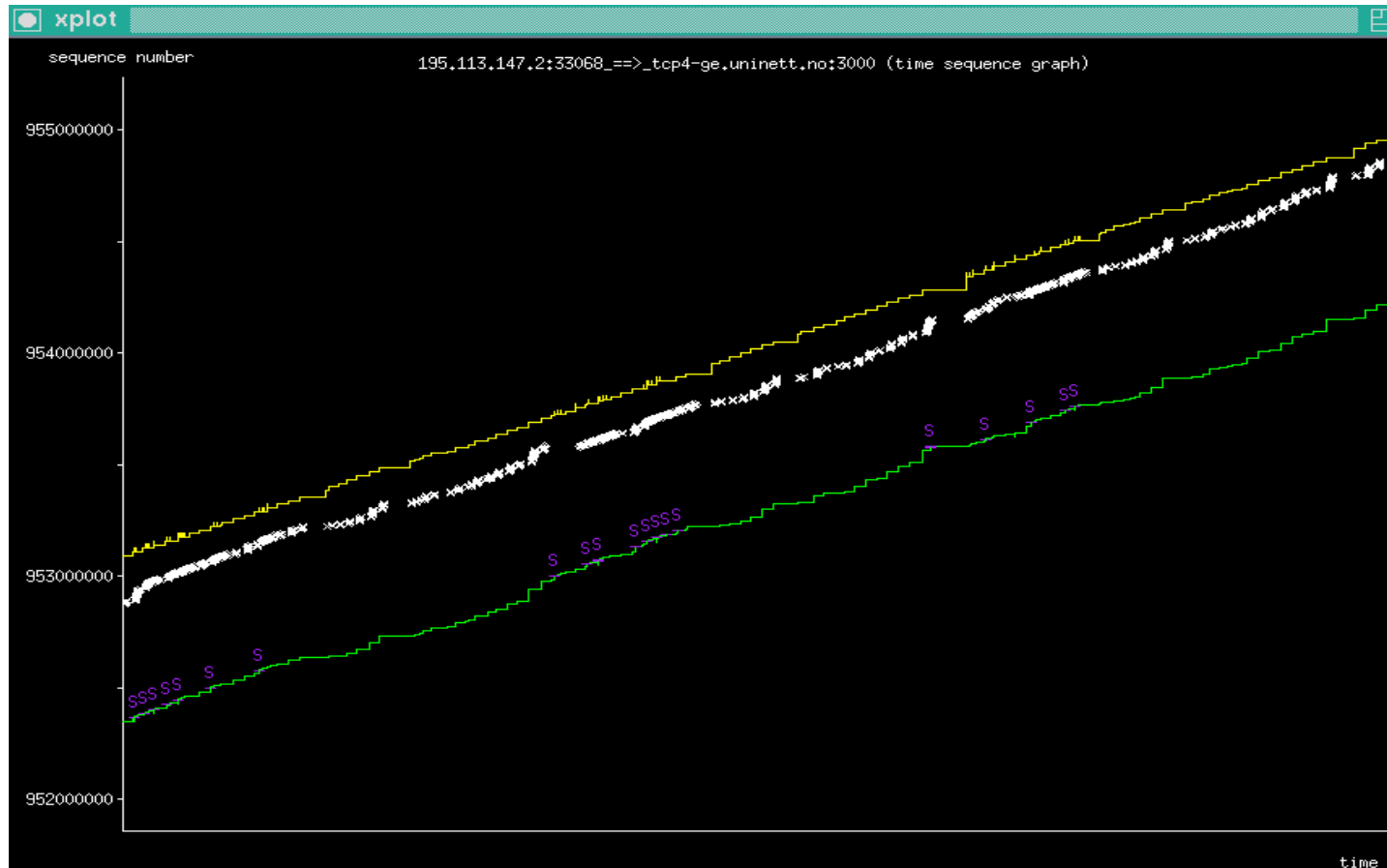
ssh performance, cont.

Bw=1 Gb/s, RTT=45 ms, TCP window=8 MB, Xeon 2.4 GHz



ssh performance, cont.

CHAN_SES_WINDOW_DEFAULT=40 * 32 kB blocks, 85% CPU load



sequence number development - rwin utilised

Conclusion

- Large TCP windows require other configuration tasks for good performance
- Buffer autoconfiguration should not just conserve memory and set TCP windows „large enough“, but also „small enough“
- ssh has a significant performance problem that must be resolved at the application level
- Influence of OS configuration and implementation-specific features on performance can be stronger than amendments in congestion control

Thank you for your attention <http://staff.cesnet.cz/~ubik>