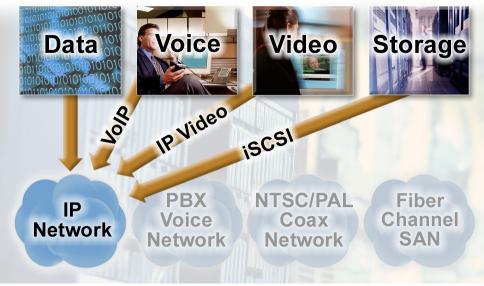


Converging Networks

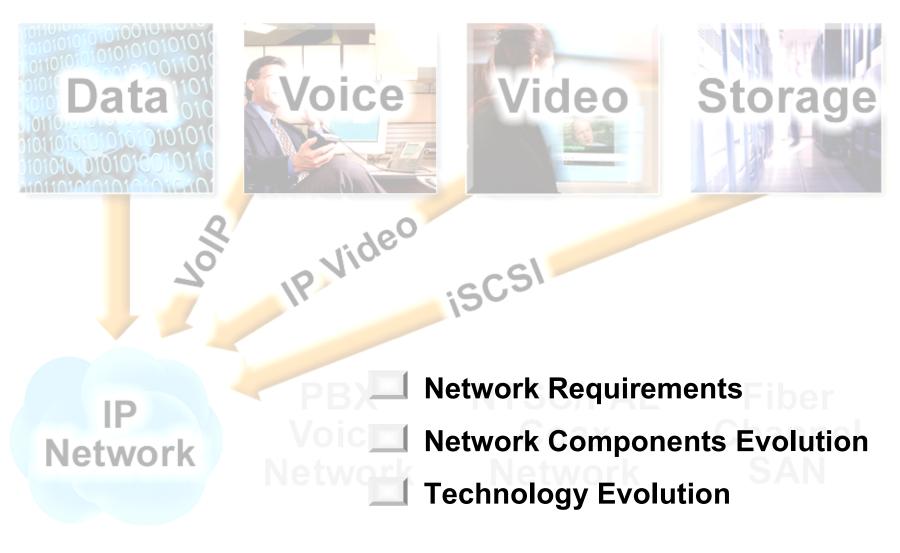


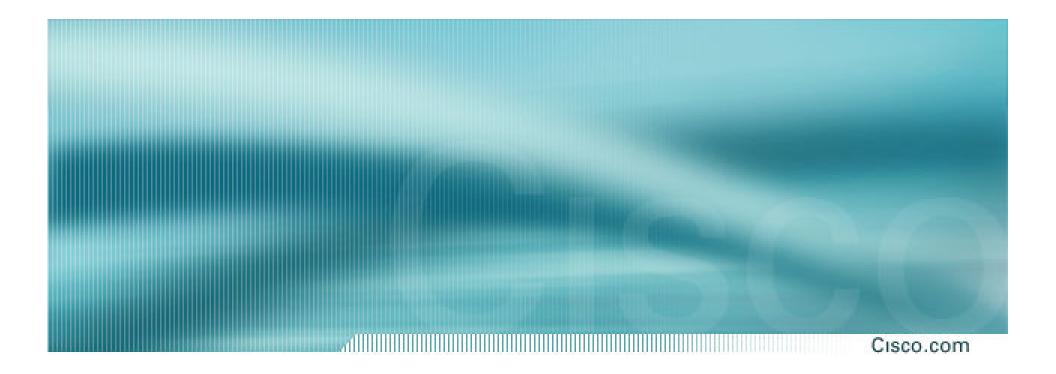
From a Technology View

Günter Honisch Distinguished Engineer EMEA Consulting ghonisch@cisco.com

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Agenda





Network Requirements

Technology Drivers

Converged networks

Voice, Video, Data, Storage Driving requirements for greater intelligence & scaleability

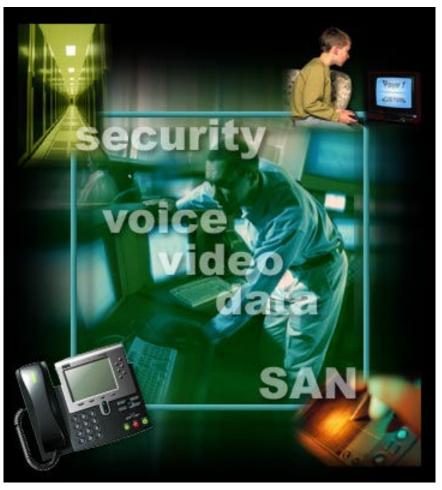
 Desktop & Server computing power increasing

GE, 10GE fullspeed connections

 Advanced applications with challenging traffic patterns

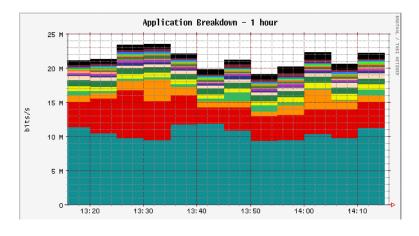
Greater reliance on QoS to manage increased bandwidth

Mission-critical security
Networks are more open
Threats ever-increasing



Network Traffic – what is going on ?

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>90% of sessions have <10 packets

Transaction mode (mail, small web page)

Stresses flowbased mechanisms in the network

>70% of all TCP traffic results from <10% of the sessions, in high rate bursts

High speed aggressive flows need solid behaviour of forwarding infrastructure (Buffers, Tail drop vs xRED)

TCP & UDP at high speeds are challenging to integrate

Applications Today

What is important:

Varies by network and population

Key applications:

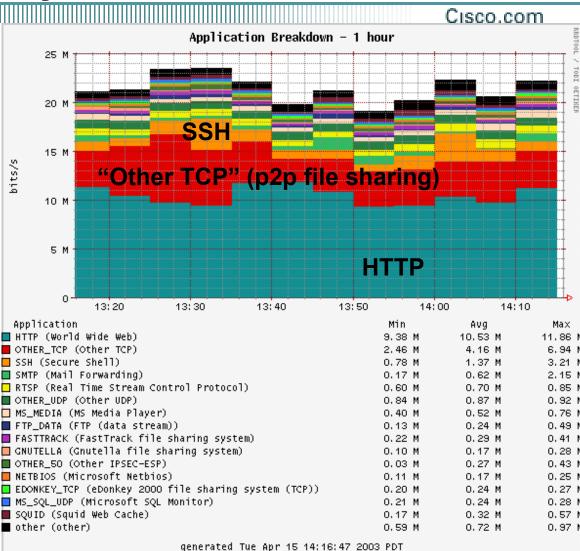
The Web+Mail

- Large file transfers
- Interactive access

Streaming audio/video

ERP

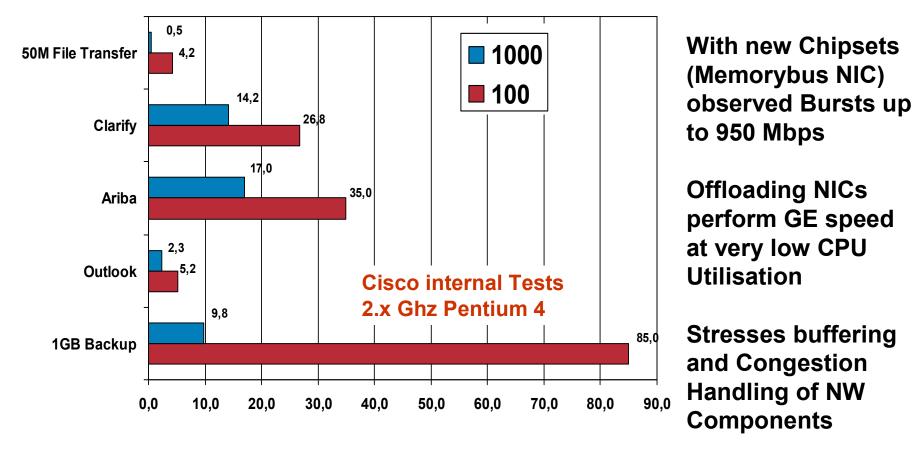
Games



http://www.caida.org/dynamic/analysis/workload/sdnap/0_0_/ts_top_n_app_bytes.html

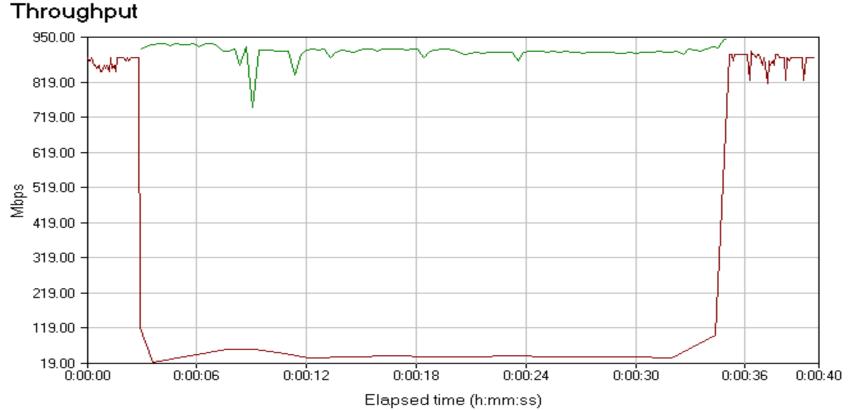
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GE on the Desktop ?



Time in Seconds

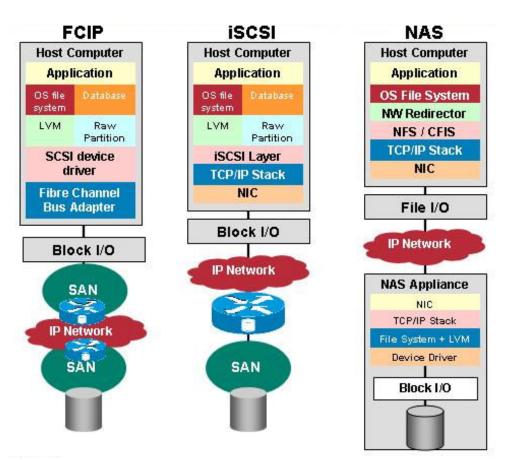
Gigabit and NO QoS...



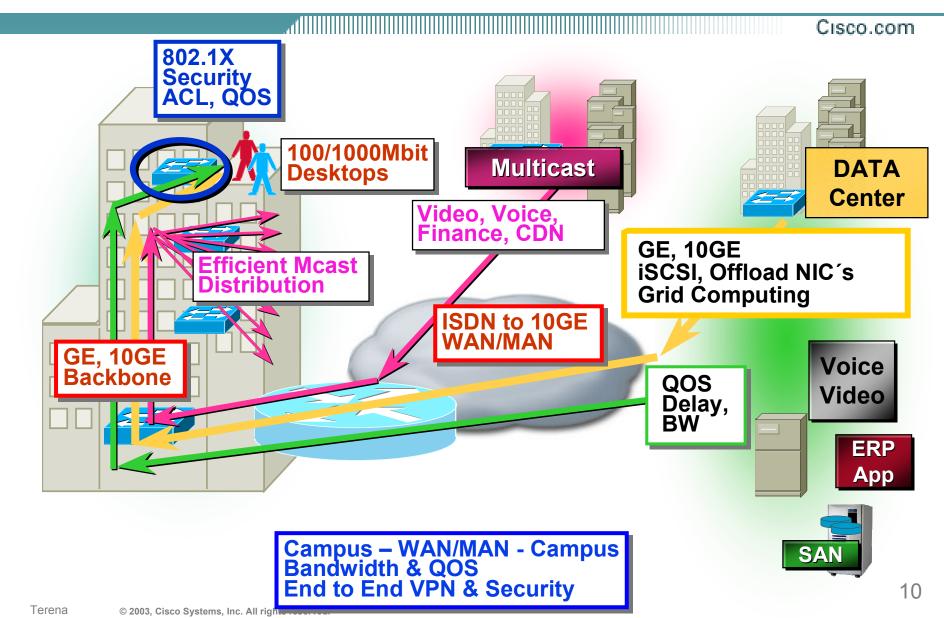
- TCP (red) and UDP (green) Streams with Gig attached hosts
- 1GB data using TCP with 0MB Loss
- 3.7GB data using UDP with 23MB lost
- 15K of 22.5M datagrams lost 154 max consecutive loss

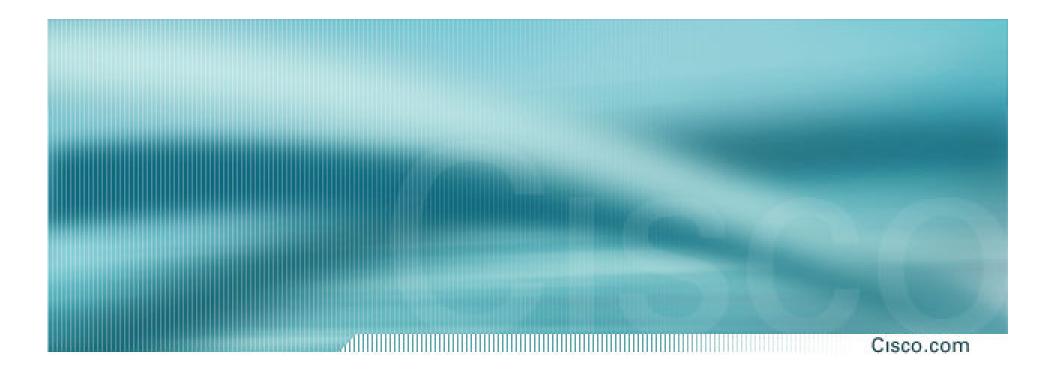
Storage over IP

- Ubiquitous: Access storage from campus, MAN, and WAN
- Agnostic: Single access technology for block and file
- Scalable: Distance, node count and performance
- Economical: Utilize IP / GE infrastructure and expertise
- Traffic: very high Bandwidth up to ~ 10Gbps in the future
- **QOS:** ideally sub Millisecond...



Increasing Requirements

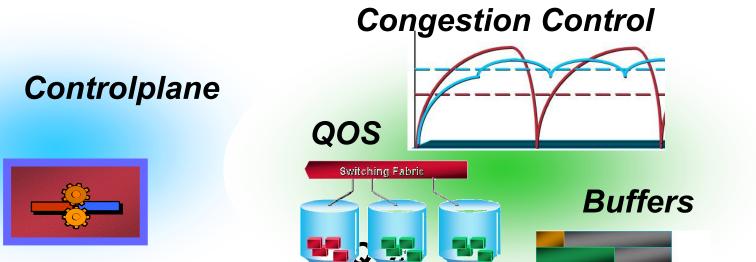


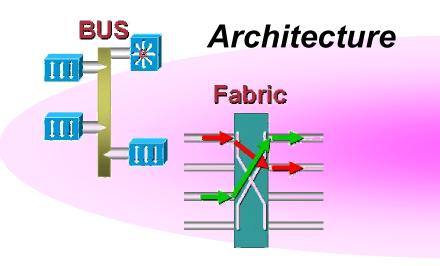


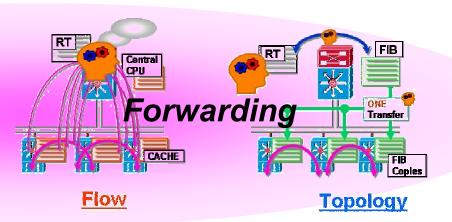
Network Component Evolution

Building Network Componentssome Decisions...



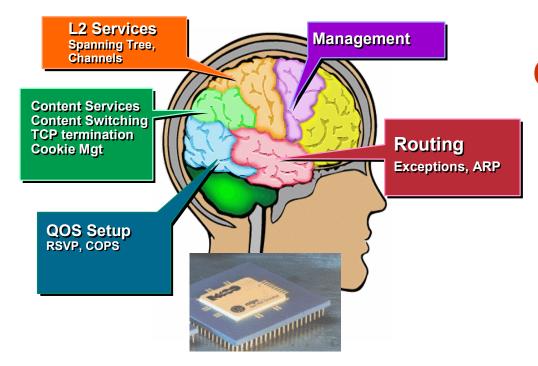






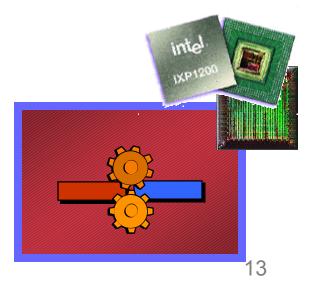
Dataplane

It is all about Brains... Controlplane - Dataplane

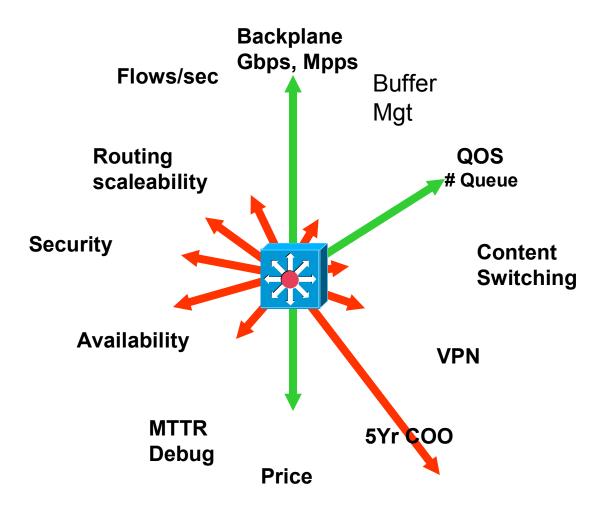


Control-Plane "SW" RISC/CISC CPU based Routing Spanningtree Housekeeping Management

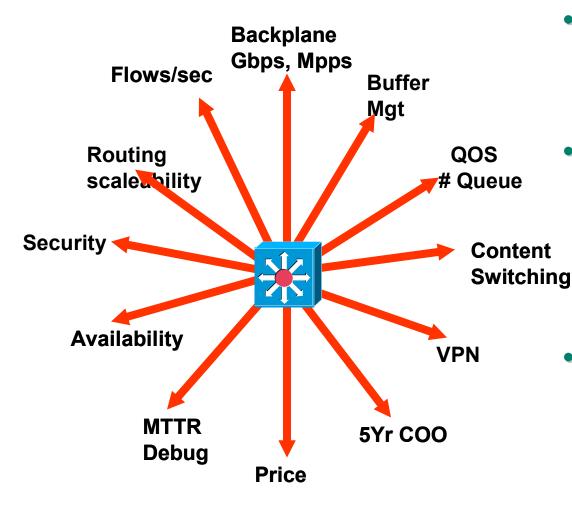
Dataplane "HW" Switching in ASICS or NW Proc forwards & manipulates Data handles Access Control, QOS



True Scaleability ?



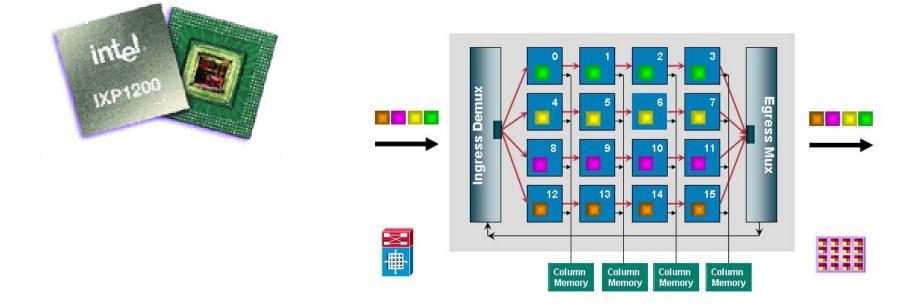
True Scaleability !



- Scaleability deals with all networking layers
- It is important to understand the different components of End to End Performance
- Focus on just one of the many variables does not help in real life situations

Network Processor

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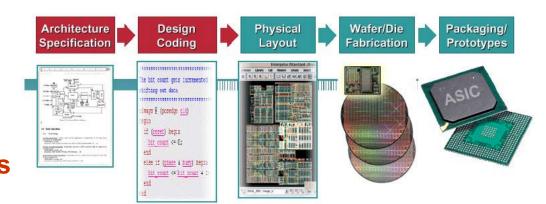
Shorter Development Cycles Changes possible More difficult to predict performance

ASIC Technology

Catalyst FFE Technology = 0.18u Die size = 13.4 x 13.4 Transistors = 91M Pin Count: 1188



Long Development Cycles Difficult to change Very high Performance Costefficient in high numbers

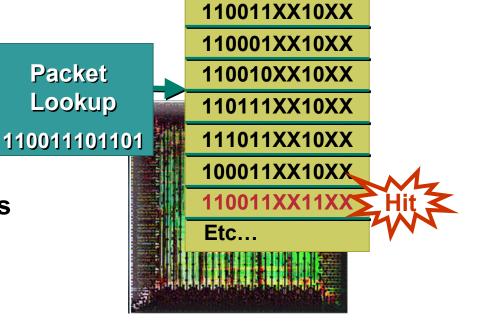


Efficient Lookup – Ternary CAM

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- Used for Routing Decisions, Access Control Lists, QOS, L2 Tables
- All entries are checked in parallel Same performance independent of number of entries
- Advantages

Longest match lookup One lookup—Fixed latency flexibility to ignore fields

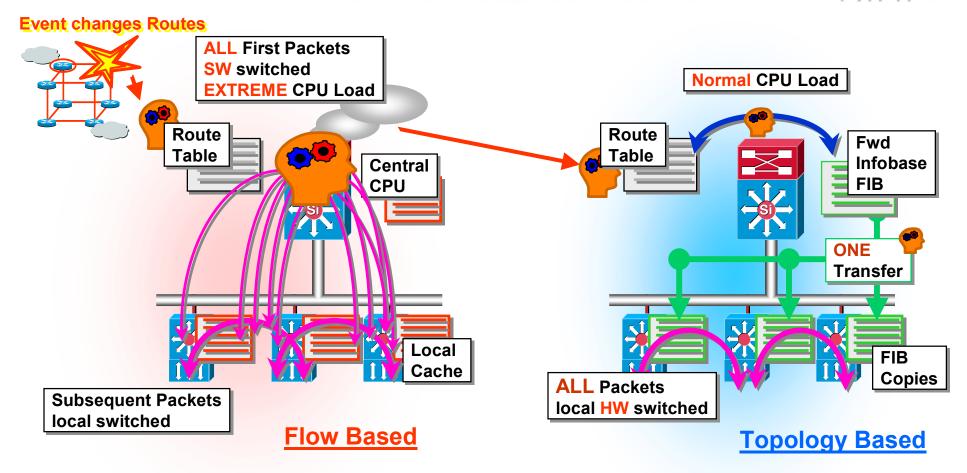


TCAM HW CEF Implementation

IP DA 10.1.1.102 S/D IP 0 10.1.1.100 /32 . L4 Skt 10.1.1.101 /32 **D-IP Addr** MASK TTL Load MAC S/D 10.1.1.0 124 **Balance** Adj Ptr VLAN # Hash # MASK Encap Result ACL 10.1.1.0 /16 . MASK . **FlowTable Adjacancy** CEF **Table** RAM TCAM MAX

SCALEABLE Forwarding why CPU Flow based mechanisms do not scale

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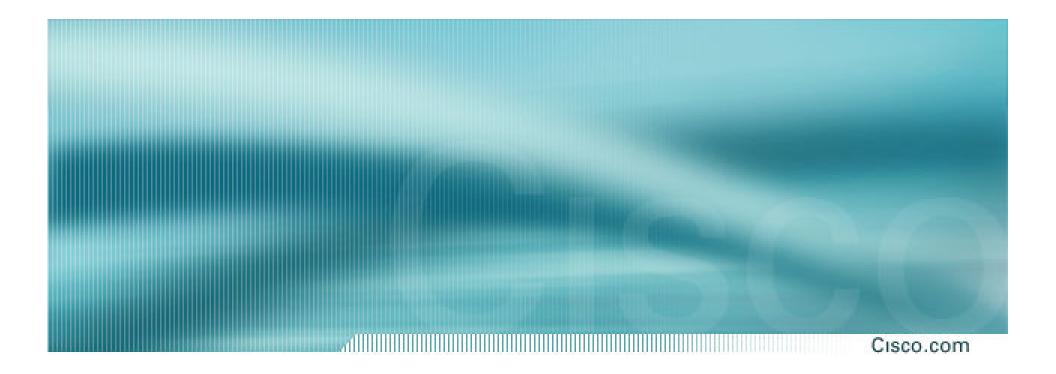


Distributed FLOW based Forwarding will never scale to Real World Traffic Patterns I.e. 100 Mio PPS would require 2 MPPS CPU performance for Flows with average 50 Pkts MOST Flows are much shorter - ALL CPUs are slower today...

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The Weak Link in the Chain Flow Based Models

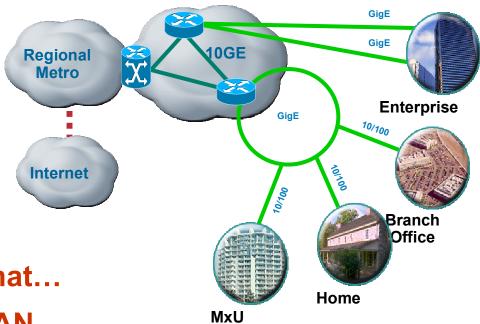
- The "First" packet in a flow must be presented to the CPU where a forwarding decision is made & subsequent packets are handled in hardware.
- CPU performance is small compared to ASIC system.
- Problems dealing with Network Reconfigurations due to building Route Tables, all new flows & cache purge all at the same time
- Many new flows may overwhelm processor, causing packet loss of data as well as control plane traffic (such as routing protocols)
- To ensure data integrity, a certain amount of maintenance "cache churn" is expected to keep the state of the flow cache current. Reflected as CPU utilization.



Technology Evolution

Ethernet Is Extending Reach

- Price/Performance
- Consistency
- Services ubiquity
- Perceived Simplicity
- BUT
- Ethernet is just a Frame Format...
- SP-Class Ethernet in WAN/MAN is different than simple Campus
- Fundamental Issues like RTT size Buffers, WAN Class QOS still apply



What is Quality of Service?

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The Pragmatic Answer: QoS is managed fairness of resources

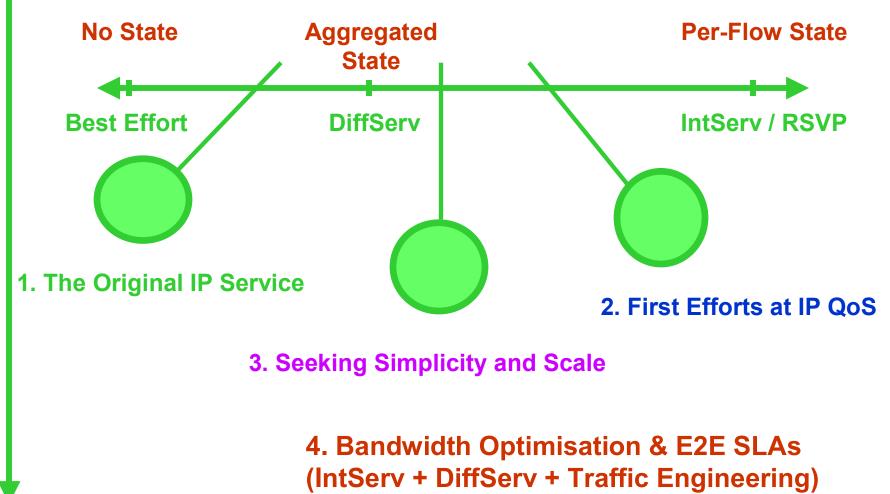
The Technical Answer:

Set of techniques to manage Delay, Jitter, Packet Loss, and Bandwidth for Flows in a Network

The IP QoS Pendulum

41111111

Time



Tight SLA – QOS

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VOIP Class SLA

Max latency 15 ms and no drops when rate <= policy bandwidth

Business Latency-Sensitive Class (EF) SLA

- latency below 30ms when no overload
- If other classes idle, take 100% of the bandwidth
- Maximum latency when this class is overloaded: < 30ms

Business Throughput-Sensitive Class (AF) SLA

minimum bandwidth guarantee when link is congested

Default Class (Best Effort) SLA

• quantum and bandwidth at remaining percentage

Tight SLA – Fast Convergence

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- Eliminate transient loops and black holes
- High availability requirements

99.999% per day ⇔ 0.9 sec of downtime

VoIP requirements

40msec Loss of Connectivity : glitch

1-2 sec Loss of Connectivity : call drop

Improvements

SPF (OSPF, IS-IS) Optimisation

IP Event Dampening

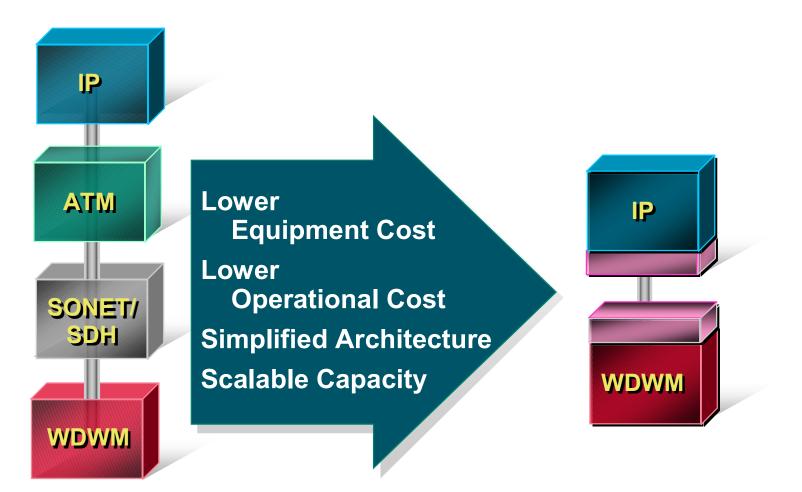
Multicast Sub-Second Convergence

BGP Optimisation

MPLS Fast Re-Route for Sub-100ms Restoration

In the Beginning: Optical OXC promised elimination of Layers

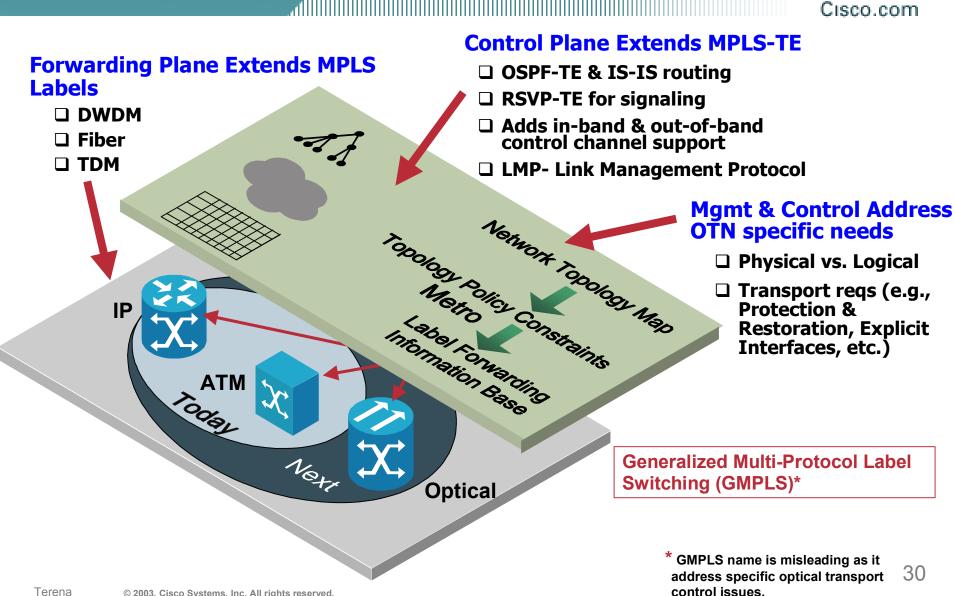
dillining Cisco.com



The Promise of a Next Gen Optical Control plane:

- Protocols to control "optical" transport networks
 - Leverage new advancements in optical technology
 - Data optimized architectures
 - handle unpredictable data traffic
 - economical for low-revenue services
 - Future-proofed, open architectures
- Protocols that better integrate IP and Optical
 - Integrated control across data and transport
 - New differentiated service and business models
 - Simplified, automated operational processes & systems

Extending MPLS Protocols for the Optical/Unified Control Plane



Ongoing Challenges to UCP Evolution

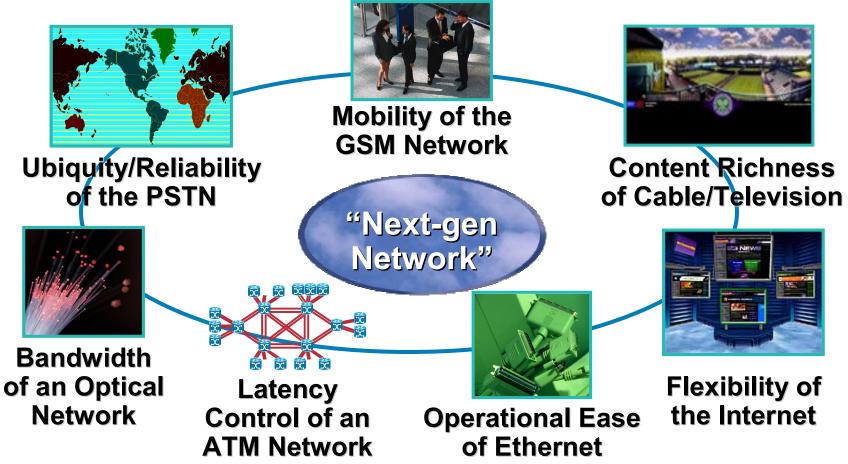
2. Telecom Bankruptcies &/or Business **Excess Bandwidth** 1. **Restructuring = Disruption** in Network Core 2. networks WORLDCOM **Global Crossing**[®] E b o n e concert The experienced broadband service liams 3. Lack of Capital, 4. Focus on near term 5. Concerns w/ Vendor **Reduced R&D Budgets Revenue; Leverage** Longevity current assets 408-745-20

31

15 HJK 1.25 RTY 1.23 IOP .05 BNM 12.0 XCV .20

So What Are the Characteristics of the Ideal Converged Next-gen Network?

Fusing the Best Properties of Today's Networks onto a Common Lowest Cost Infrastructure





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